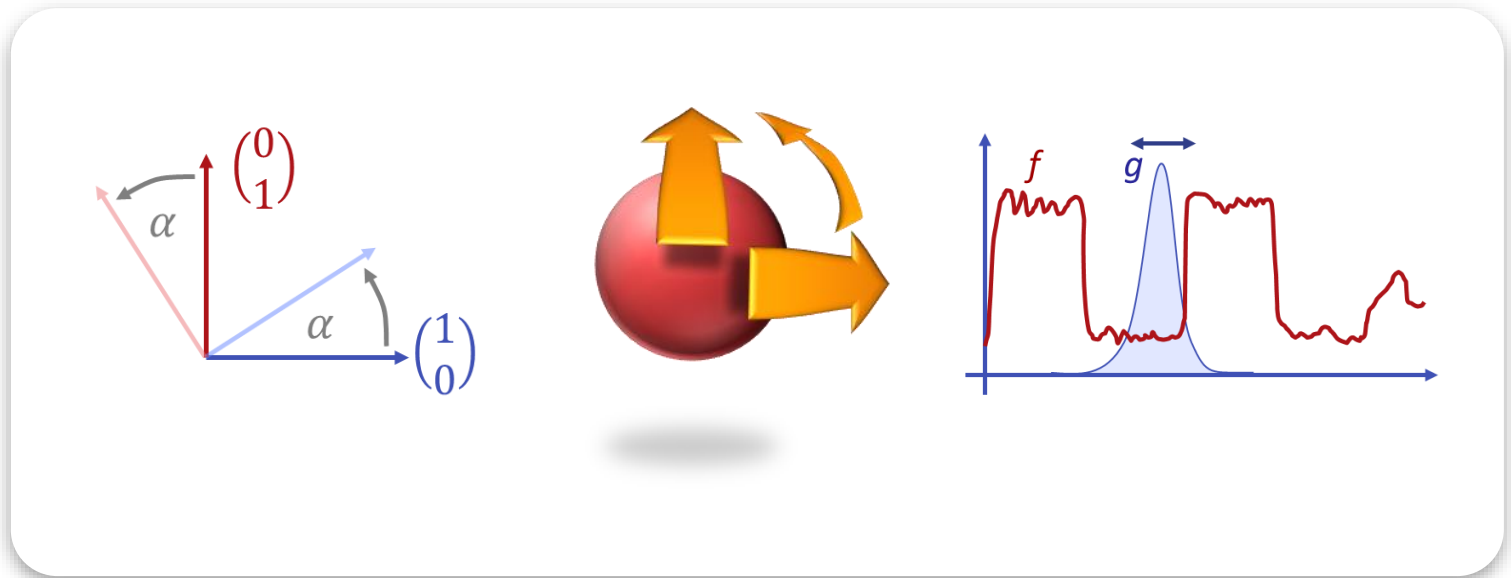


Modelling 1

SUMMER TERM 2020



LECTURE 8

(Linear) Information Loss

Information Loss in Linear Mappings

Linear Maps

A function

- $f: V \rightarrow W$ between vector spaces V, W

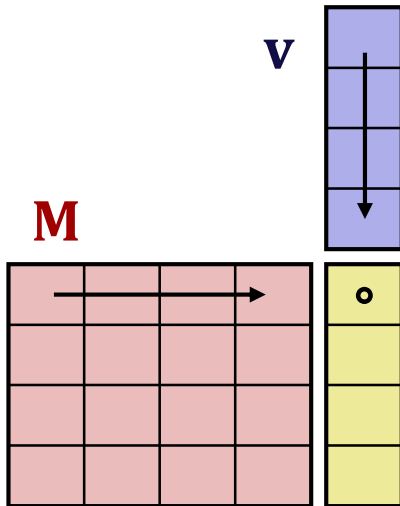
is linear if and only if:

- $\forall \mathbf{v}_1, \mathbf{v}_2 \in V: f(\mathbf{v}_1 + \mathbf{v}_2) = f(\mathbf{v}_1) + f(\mathbf{v}_2)$
- $\forall \mathbf{v} \in V, \lambda \in \mathbb{R}: f(\lambda \mathbf{v}) = \lambda f(\mathbf{v})$

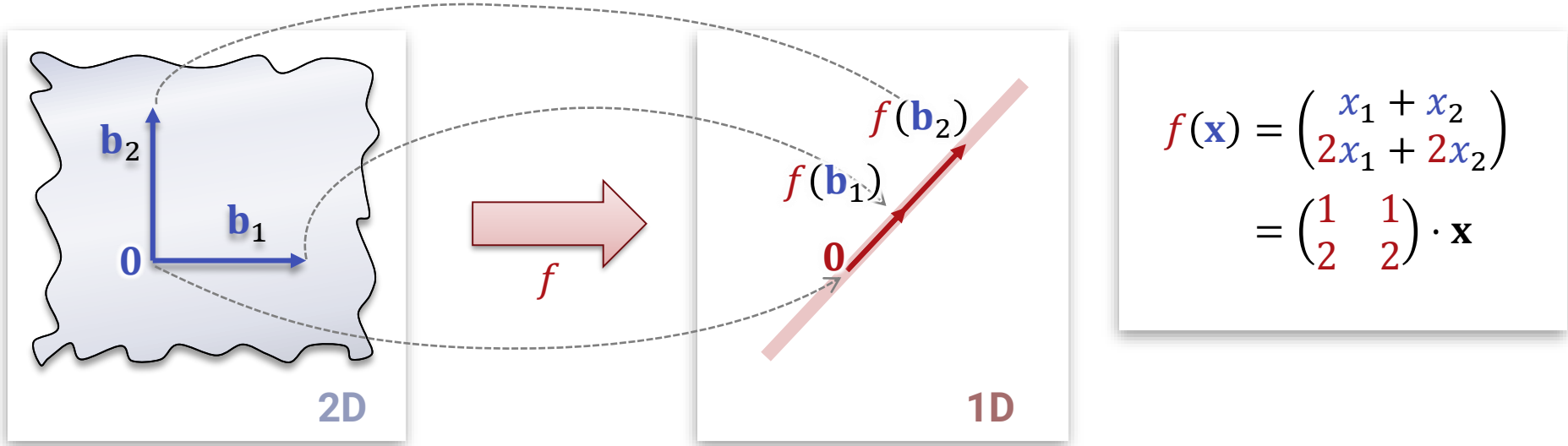
Matrix Product

All operations are matrix-matrix products:

- Matrix-Vector product:
- $f(\mathbf{x}) = \mathbf{M}_f \cdot \mathbf{x}$



Not invertible

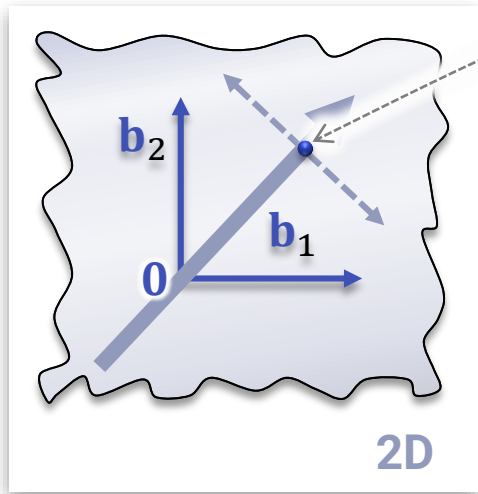


$$\begin{aligned} f(\mathbf{x}) &= \begin{pmatrix} x_1 + x_2 \\ 2x_1 + 2x_2 \end{pmatrix} \\ &= \begin{pmatrix} 1 & 1 \\ 2 & 2 \end{pmatrix} \cdot \mathbf{x} \end{aligned}$$

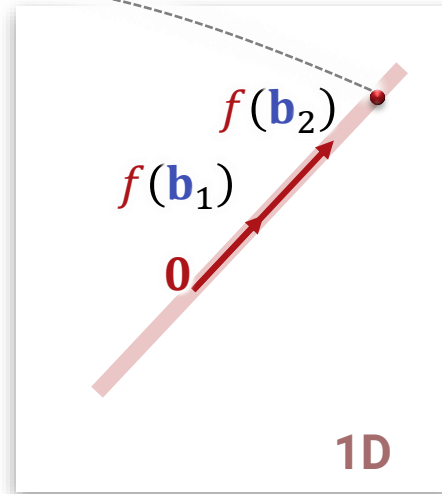
Information flow:

- After f , we can recover $b_1 + b_2$
 - Sum of inputs
- We do not know $b_1 - b_2$ anymore
 - Difference of inputs

Not invertible



f^{-1} | where possible

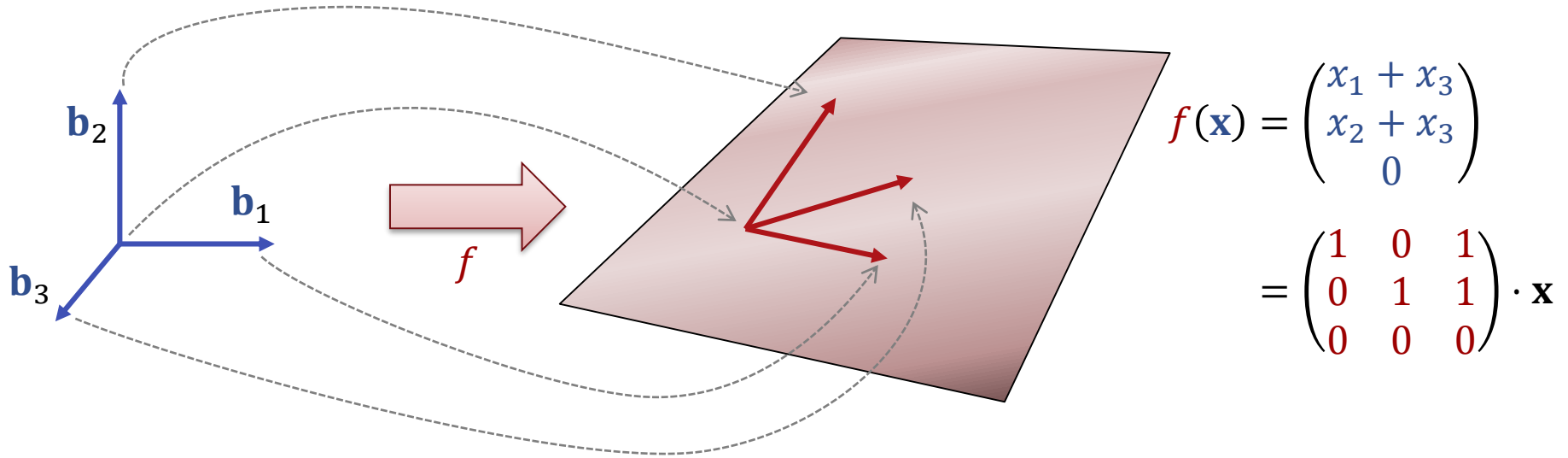


$$\begin{aligned} f(\mathbf{x}) &= \begin{pmatrix} x_1 + x_2 \\ 2x_1 + 2x_2 \end{pmatrix} \\ &= \begin{pmatrix} 1 & 1 \\ 2 & 2 \end{pmatrix} \cdot \mathbf{x} \end{aligned}$$

Information flow:

- After f , we can recover $b_1 + b_2$
 - Sum of inputs
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 - Difference of inputs

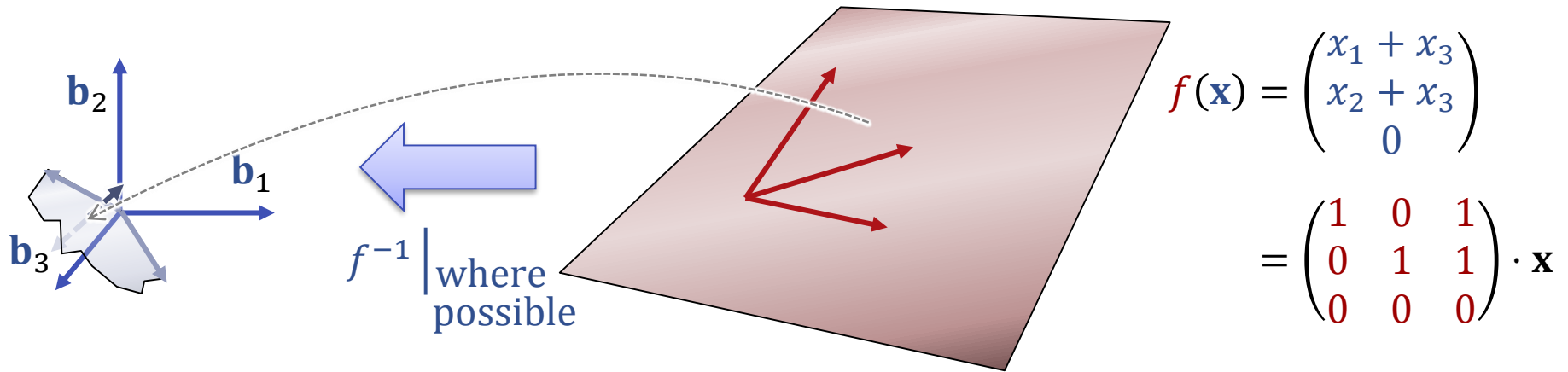
Not invertible



Information flow:

- After f , we can recover $b_1 + b_3$ and $b_2 + b_3$
- We do not know $b_2 - b_3$ anymore

Not invertible



Information flow:

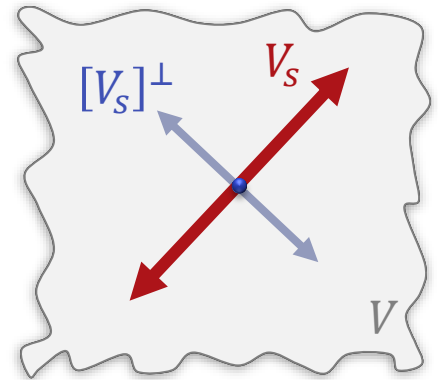
- After f , we can recover $b_1 + b_3$ and $b_2 + b_3$
- We do not know $b_2 - b_3$ anymore

Orthogonal Complement

Definition

- **Given:** Subspace $V_S \subseteq V$
- Orthogonal complement

$$V_S^\perp := \{\mathbf{v} \in V \mid \forall \mathbf{w} \in V_S: \langle \mathbf{v}, \mathbf{w} \rangle = 0\}$$



Intuition

- Set of all vectors orthogonal to V_S
- Zero projection onto any $\mathbf{w} \in V_S$

Theorem

$$V_S \subset V \Rightarrow V = \text{span}\{V_S, V_S^\perp\} \quad [:= V_S \oplus V_S^\perp]$$

In general

Consider mapping

$$f: V_1 \rightarrow V_2$$

Subspaces of V_1

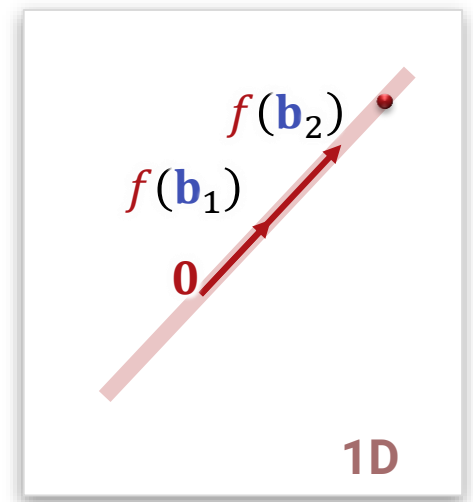
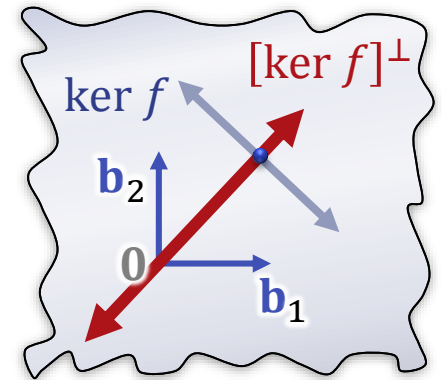
- **Kernel:** Subspace that is lost

$$\ker f := \{\mathbf{x} \in V_1 \mid f(\mathbf{x}) = 0\}$$

- **Orthogonal complement of kernel**

$$[\ker f]^\perp = \{\mathbf{v} \in V_1 \mid \forall \mathbf{w} \in \ker f: \langle \mathbf{v}, \mathbf{w} \rangle = 0\}$$

- In this space, f is invertible



In general

Consider mapping

$$f: V_1 \rightarrow V_2$$

In the target domain

$$\text{im } f := \{\mathbf{y} \in V_2 \mid \exists \mathbf{x} \in V_1: f(\mathbf{x}) = \mathbf{y}\}$$

- Subspace of V_2
- Same dimension as kernel complement

$$\dim([\ker f]^\perp) = \dim(\text{im } f)$$

In general

Consider mapping

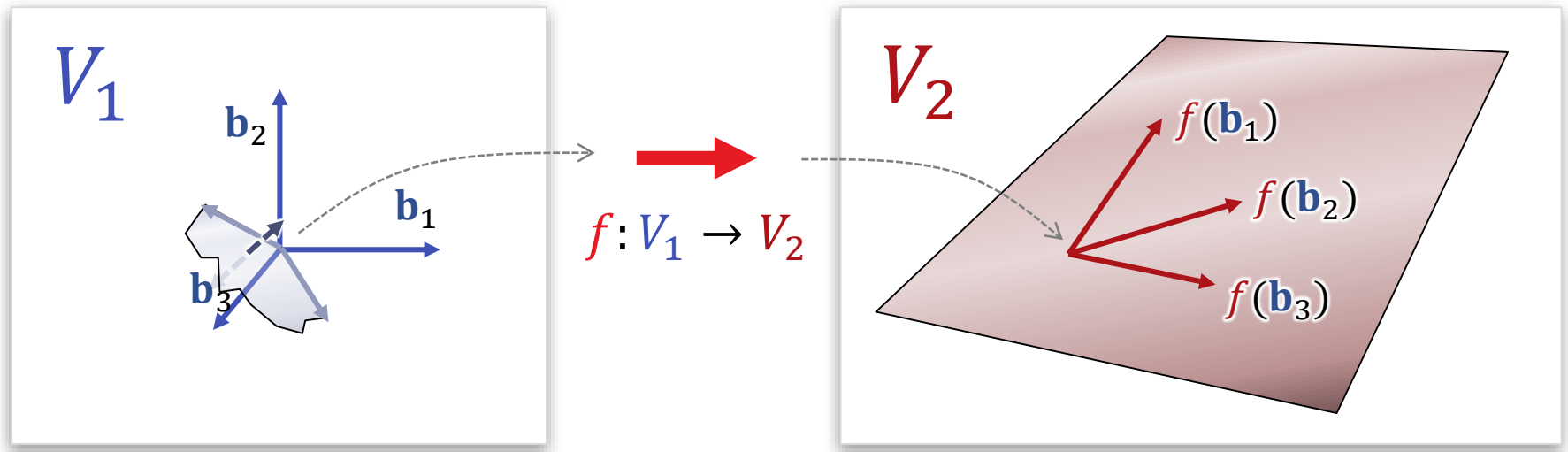
- Rank is the dimension of the mapped space

$$\begin{aligned}\text{rank}(f) &:= \dim(\text{im } f) \\ &= \dim(\text{span}(V_1 \setminus \ker f))\end{aligned}$$

- Source space V_1 is split:
 - $\dim \text{im}(f)$ = dimensions “preserved” by f
 - $\dim \ker(f)$ = dimensions “removed” by f
- Sums up:

$$\dim(V_1) = \dim(\text{im } f) + \dim(\ker f)$$

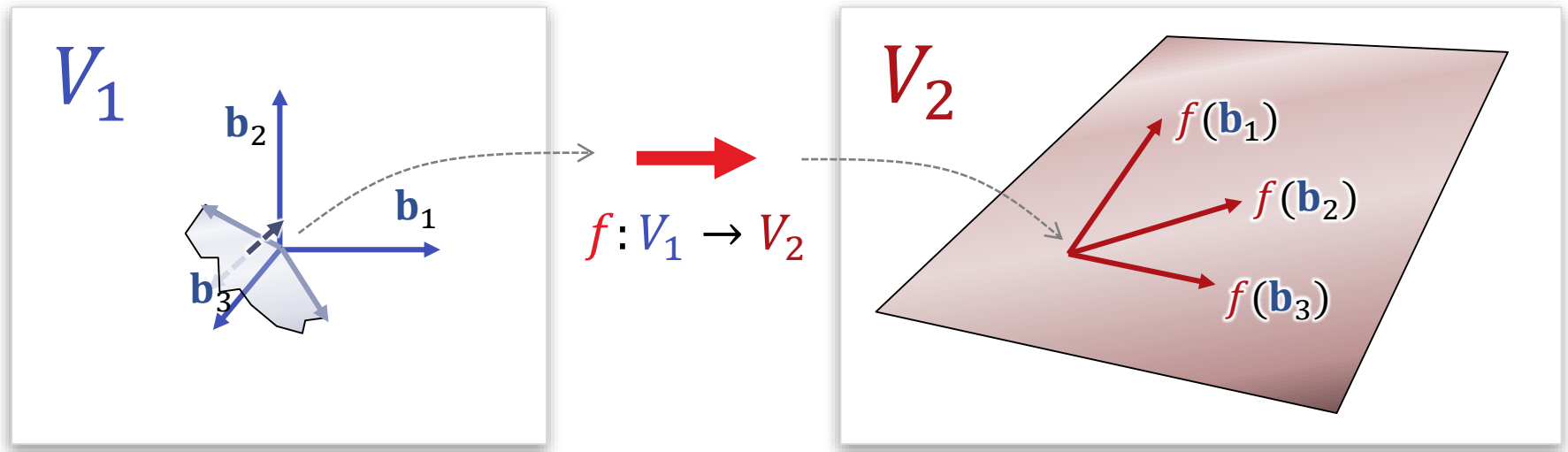
Structural Insight



Mapping Subspaces to Subspaces

- Invertible map from $[\ker f]^\perp \rightarrow \text{im } f$
- Not covered
 - “Source” information lost: coordinates within $\ker f$
 - Unreachable “targets”: vectors within $[\text{im } f]^\perp$

Structural Insight



Dimensions add up

- $\dim[\ker f]^\perp = \dim \mathbf{im} f$
- $\dim V_1 = \dim \ker f + \dim[\ker f]^\perp$
- $\dim V_2 = \dim \mathbf{im} f + \dim[\mathbf{im} f]^\perp$

In practice?

In practice

- It always never works:
 - Most matrices have noise (measurement, numerics)
 - Any practical mapping has “full rank”
 - Inverting matrices is not always stable
 - Even full-rank matrices might delete information
 - Need to understand this better!

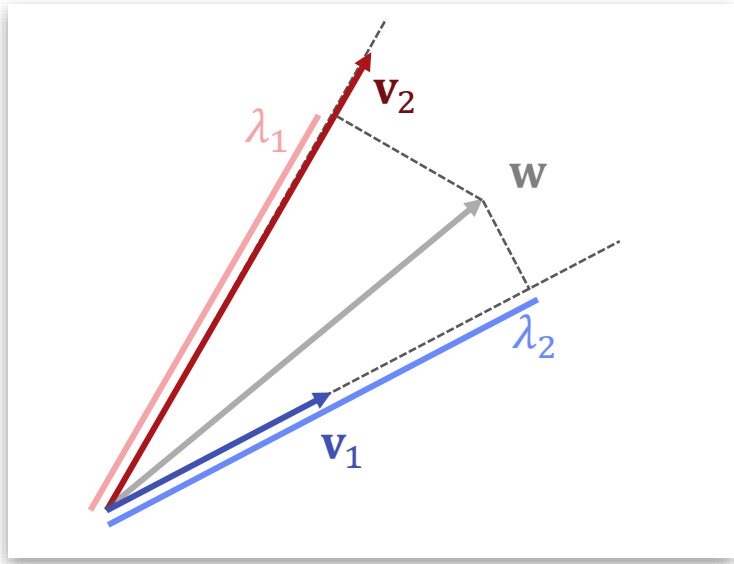
We will discuss this soon

- Tools:
 - Eigenvalues
 - Singular value decomposition (SVD)

Linear Systems of Equations
Inverting Linear Maps

Situation

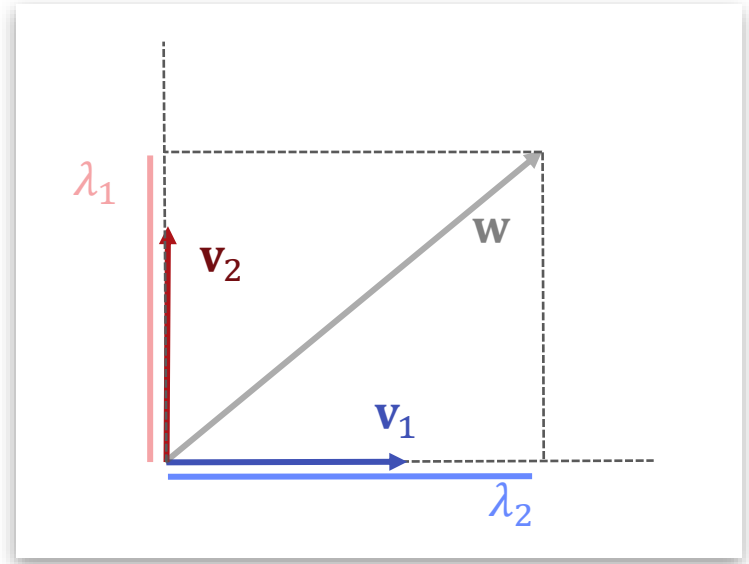
General Case



Linear System

$$\lambda_1 \cdot \mathbf{v}_1 + \dots + \lambda_n \cdot \mathbf{v}_n = \mathbf{w}$$

Orthogonal



Direct Computation

$$\lambda_1 = \mathbf{v}_1 \cdot \mathbf{w}$$

\vdots

$$\lambda_n = \mathbf{v}_n \cdot \mathbf{w}$$

Linear Systems of Equations

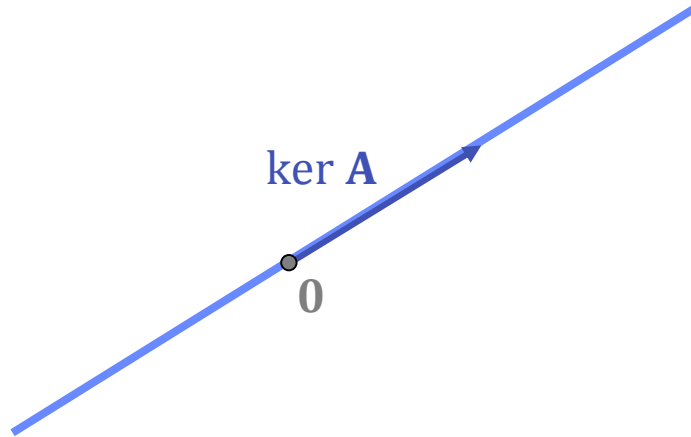
Problem: Invert an affine map

- Given: $\mathbf{A} \cdot \mathbf{x} = \mathbf{b}$, i.e., $\mathbf{A} \cdot \mathbf{x} - \mathbf{b} = \mathbf{0}$
 - We know \mathbf{A}, \mathbf{b}
 - Looking for \mathbf{x}
- Compute $\mathbf{x} = \mathbf{A}^{-1} \cdot \mathbf{b}$

Solution

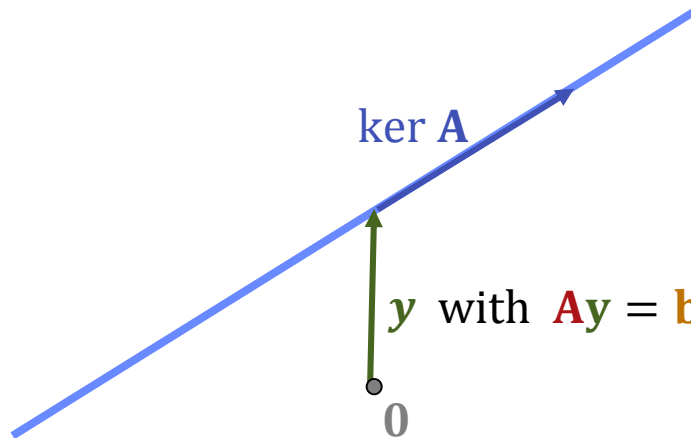
- Set of solutions: *affine subspace* of \mathbb{R}^n (or \emptyset)
 - Point, line, plane, hyperplane...
- Innumerable algorithms

Linear Systems of Equations



$\{\mathbf{x} \mid \mathbf{Ax} = \mathbf{0}\}$ – hyperplane through the origin

“Homogeneous” system



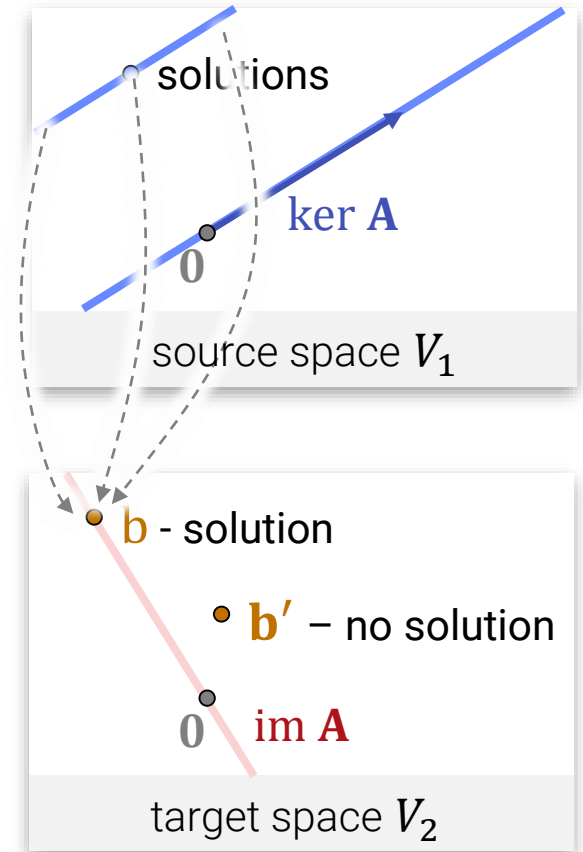
$\{\mathbf{x} \mid \mathbf{Ax} = \mathbf{b}\}$ – hyperplane through any point

“Inhomogeneous” system

Structure

Linear System ($\mathbf{A}: V_1 \rightarrow V_2$):

- $\mathbf{Ax} = \mathbf{0}$
 - Solution space = $\ker \mathbf{A}$
- $\mathbf{Ax} = \mathbf{b}$
 - Might or might not have a solution
 - Solution if and only if $\mathbf{b} \in \text{im } \mathbf{A}$
- Set of all solutions:
 - One \mathbf{y} with $\mathbf{Ay} = \mathbf{b}$
 - Add any solution of $\mathbf{Ax} = \mathbf{0}$
 - Solution set: $\mathbf{y} + \ker \mathbf{A}$

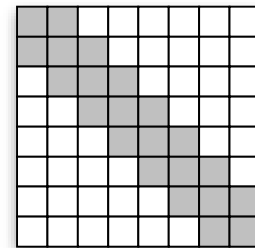


Solvers for Linear Systems

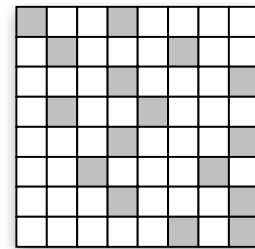
Solving linear systems of equations

- **Baseline:** Gaussian elimination
 $O(n^3)$ operations for $n \times n$ matrices
- We can do better, in particular for special cases:

- **Band matrices:**
constant bandwidth



- **Sparse matrices:**
constant number of non-zero entries per row
 - Store only non-zero entries



Solvers for Linear Systems

Algorithms: linear systems of n equations

- Band matrices, $O(1)$ bandwidth:
 - Modified $O(n)$ elimination algorithm.
- Iterative Gauss-Seidel solver
 - converges for diagonally dominant matrices
 - Typically: $O(n)$ iterations, each costs $O(n)$ for a sparse matrix.
- Conjugate Gradient solver
 - Only symmetric, positive definite matrices
 - Guaranteed: $O(n)$ iterations
 - Typically good solution after $O(\sqrt{n})$ iterations.

See: [J. R. Shewchuk, An Introduction to the Conjugate Gradient Method Without the Agonizing Pain, 1994.](#)